#### Overview

- Last Lecture
  - Data Integrity 1
- This Lecture
  - Data Integrity 2
  - Source: Sections 10.2, 10.5
- Next Lecture
  - Data Security 1
  - Source: Sections 31.1, 31.2

#### Last Lecture: Parity Checking



## Simple Parity Checking (Last Lecture)

• Simple parity checking: Add an extra bit (parity bit) to the data, such that the total number of 1 bits is even (even parity) or odd (odd parity)

Even parity		Odd Parity	
1101 0100	0	1101 0100	1

- Limitations of Simple Parity Checking
  - Cannot detect an even number of bit errors
  - Burst errors (Many bits are damaged)
  - Does this make it useless?

## Parity Checking (Double Bit Error)

• Double Bit Error Detection: Use two parity bits (One for even numbered bits, One for odd numbered bits). If one bit errors or some double bit errors occur, they are detected





## Parity Checking (Multiple Bit Errors)

- What happens if two even numbered (or odd numbered) bits are in error?
- Multiple bit errors: A two-dimensional bit array
  - Rows byte or several bytes
  - Parity bit for each row
  - Transmit the bits by column
- Can detect any single burst error whose duration is less than the time to send a column.
- Parity Checking forms a basis for an error correction scheme

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#### Parity Checking (Multiple Bit Errors, cont.)

#### Sender Receiver Row (frame) Parity bit for Row Parity bit number one row number for one row 1\* Burst error occurs and Ø destroys column four, 0\* making it all zeroes. 0\* 1\* 1\* Column number Column number

#### **Figure 4.3** Detecting Burst Errors Using Parity Bits

\* Parity bit is not correct

#### Error Correction

- Can be handled two ways:
  - Ask the sender to retransmit the data
  - Correct the errors using an error correction code
- Hamming codes
  - Even parity is used in the examples
  - Assume we are sending bytes
  - Assume only one error
  - Can be extended to multiple bit errors

Hamming code for data correction is one of the most difficult parts in 244 lectures



But,



#### You are so clever to make it easy

Binary Card game: http://gwydir.demon.co.uk/jo/numbers/binary/cards.htm



#### Example

- We want to send 0110 0111
- Suppose we receive 0101 0110 0111. Is there an error? If so, what is the correct bit pattern?
- Using the Hamming code, we can not only detect single bit error, but also correct them!

#### Hamming code

- Invented by Richard Hamming in 1950
- can detect and correct on bit error



- What we will learn for Hamming code in this lecture are:
- construct Hamming code
   (e.g. data: 01100111, hamming code: 010111010111)
- detect and correct one-bit error (e.g. data received: 010101010111)

## Hamming Code

- Basic idea:
  - Use multiple extra parity bits
  - Combine bits in a smart way that each extra bit checks whether there is an error in a subset of data bits



Lecture 6 - Data Integrity 2

• A four-step approach to construct Hamming code Step 1: determine the positions of parity bits

The bits whose position number is a power of 2(1, 2, 4, 8, ..) are reserved for parity bits.



• A four-step approach to construct Hamming code Step 2: insert data bits

Insert data bits into unoccupied positions from left to right



• A four-step approach to construct Hamming code Step 3: Group the bits into different subsets with each subset containing one parity bit

Start from the  $2^{k-1}$  bit, choose  $2^{k-1}$  bits, skip  $2^{k-1}$  bits, choose  $2^{k-1}$  bits ...



• A four-step approach to construct Hamming code Step 4: calculate the values for the parity bits

Within each group, apply single parity check



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### Data to send: 0 1 1 0 0 1 1 1

# Hamming code: 0 1 0 1 1 1 0 1 0 1 1 1

#### We got the Hamming code!

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Why should Hamming code be constructed in such a way? magic trick (guess your birthday month)

























One error can be automatically corrected without the need of retransmission!



If you have n parity bits, it can identify 2<sup>n</sup>-1 bit positions

#### Hamming Code (summary)



#### Hamming Code (summary)

#### Figure 4.10 Bit Stream Before Transmission

Data:	0	1	1	0	0	1	1	1
	$m_1$	$m_2$	$m_3$	$m_4$	$m_5$	$m_6$	$m_7$	$m_8$

Hamming code:	0	1	0	1	1	1	0	1	0	1	1	1
	$p_1$	$p_2$	$m_1$	$p_3$	<i>m</i> <sub>2</sub>	<i>m</i> <sub>3</sub>	$m_4$	$p_4$	$m_5$	m <sub>6</sub>	$m_7$	$m_8$



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#### Hamming Codes (summary)

- Consider a 4 bit number b<sub>4</sub>, b<sub>3</sub>, b<sub>2</sub>, b<sub>1</sub>
- $b_i = 0$  if parity check for  $p_i$  succeeds, otherwise  $b_i = 1$

Table 4.2         Bit Position Errors and Associated Parity Error	rors
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Erroneous Bit Position	INVALID PARITY CHECKS	$b_4, b_3, b_2, \text{ and } b_1$			
No error	None				
1	$p_1$	0001			
2	$p_2$	0010			
3	$p_1$ and $p_2$	0011			
4	$p_3$	0100			
5	$p_1$ and $p_3$	0101			
6	$p_2$ and $p_3$	0110			
7	$p_1, p_2, \text{ and } p_3$	0111			
8	$p_4$	1000			
9.	$p_1$ and $p_4$	1001			
10	$p_2$ and $p_4$	1010			
11	$p_1, p_2, \text{ and } p_4$	1011			
12	$p_3$ and $p_4$	1100			

#### Hamming Code (summary)



## Multiple Bit Error Correction

- Use the idea similar to multiple bit error detection.
- Create the Hamming codes for each byte.
- Arrange the Hamming codes into an array.
- Send the array a column at a time.
- As long as no burst is longer than a column, correction can be done.

#### Error Correction vs. Error Detection

- Error detection
  - Error detection is the ability to determine that data has an error
  - Simpler, Cheaper, Requires retransmission
- Error correction
  - Error correction is the ability to correct the erroneous data.
  - No retransmission, Complex mechanism, Sometimes required

#### Summary (this lecture)

- Hamming code
  - Code construction
  - Error correction
- What should you learn from this lecture:
  - construct the Hamming code
    (e.g. data: 01100111, hamming code: 010111010111)
  - detect and correct one-bit error (e.g. data received: 010101010111)

## Summary (Data integrity)

- Error detection schemes
  - Parity checking (single bit, double bit, burst)
  - CRC (burst)
- Error correction schemes
  - Hamming codes (single bit)
- Comparison of error detection and correction