

# **COSC 301**

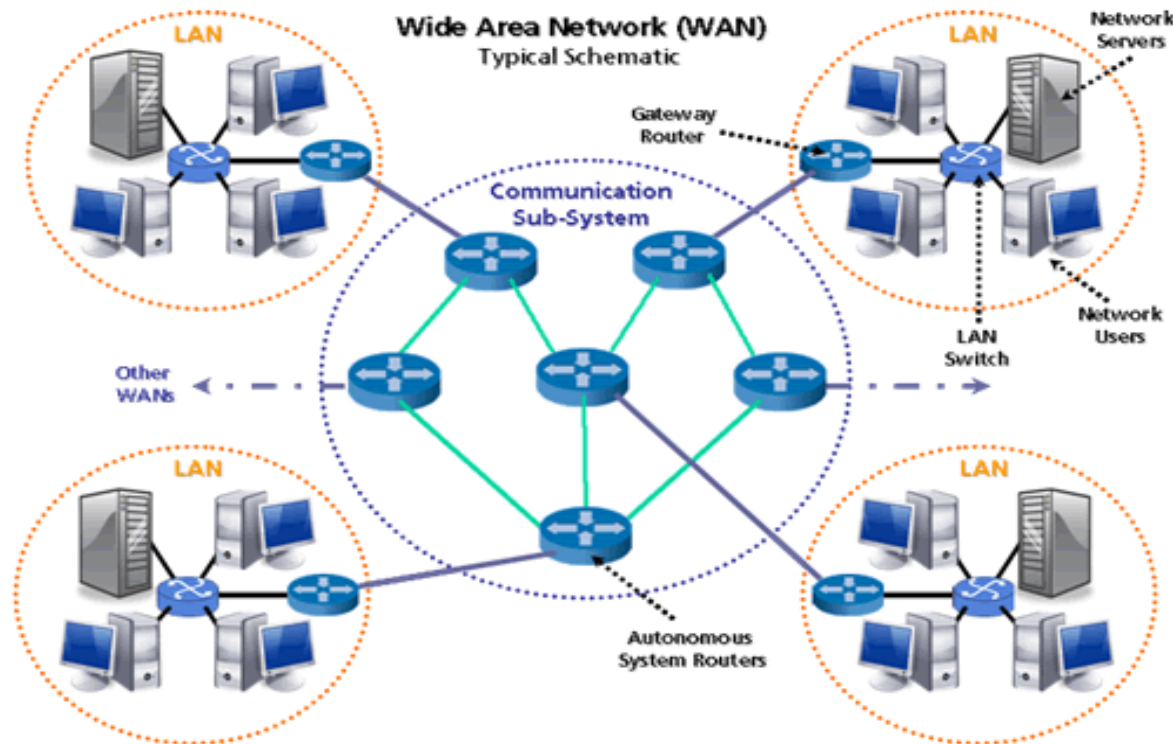
## **Network Management**

### Lecture 19: Interior Routing

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# Today's Focus



## Interior Routing

- Autonomous System
- RIP
- OSPF

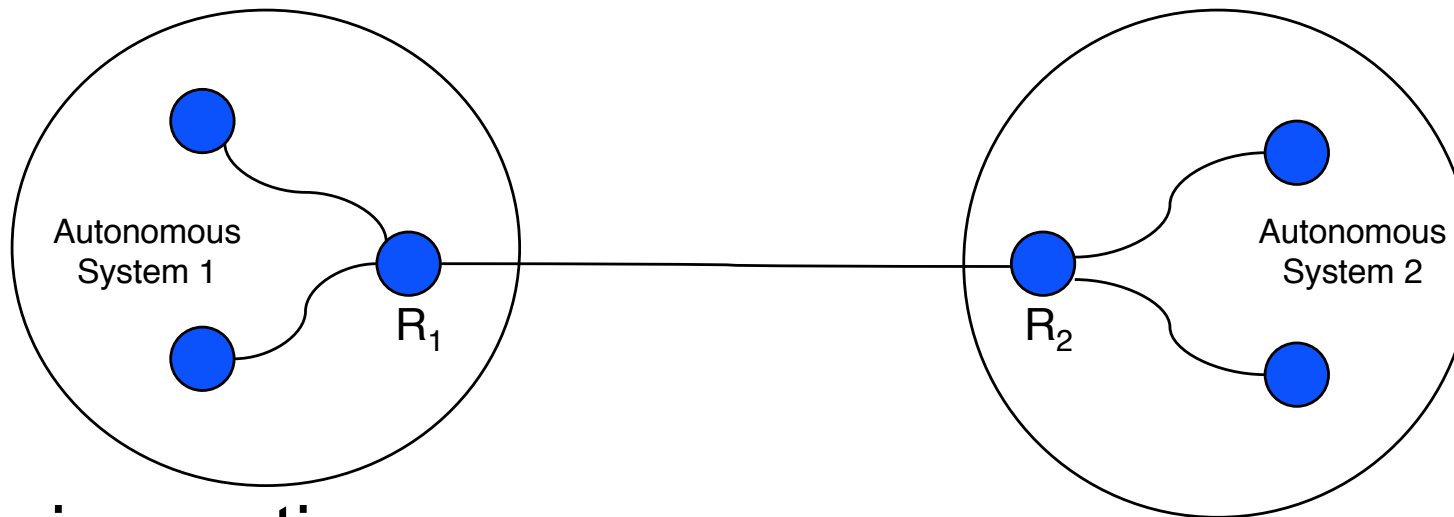
# Autonomous System

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- The Internet does not consist of independent networks
  - Networks and routers are owned by organizations or individuals
  - Networks and routers owned by a given entity fall under a single administrative authority, which guarantees that internal routes remain consistent and viable.
- Autonomous System: a group of networks and routers controlled by a single administrative authority
  - Routers are free to choose their own mechanisms for discovering, validating, and checking consistency of routes
  - The Internet is divided into autonomous systems with each owned by a single administrative authority.
  - In the current Internet, each large ISP is an autonomous system.

# Internet Routing

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- Interior routing
  - Used within an autonomous system
  - Used within an area of administrative control
  - RIP, OSPF
- Exterior routing
  - Used between autonomous systems
  - Used to peer with networks without administrative control
  - BGP

# Static vs. Dynamic

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- Static routing
  - Advantage: Simplicity, less software that might fail
  - Disadvantage: Limits on scalability, less resilience and load-balancing, slow response to handle problems
- Dynamic routing
  - Advantage: Scalable, optimal routes, automatic fail-over and fail-back, load balancing
  - Disadvantage: route update consumes bandwidth, additional load on router CPU

# Routing Information Protocol (RIP)

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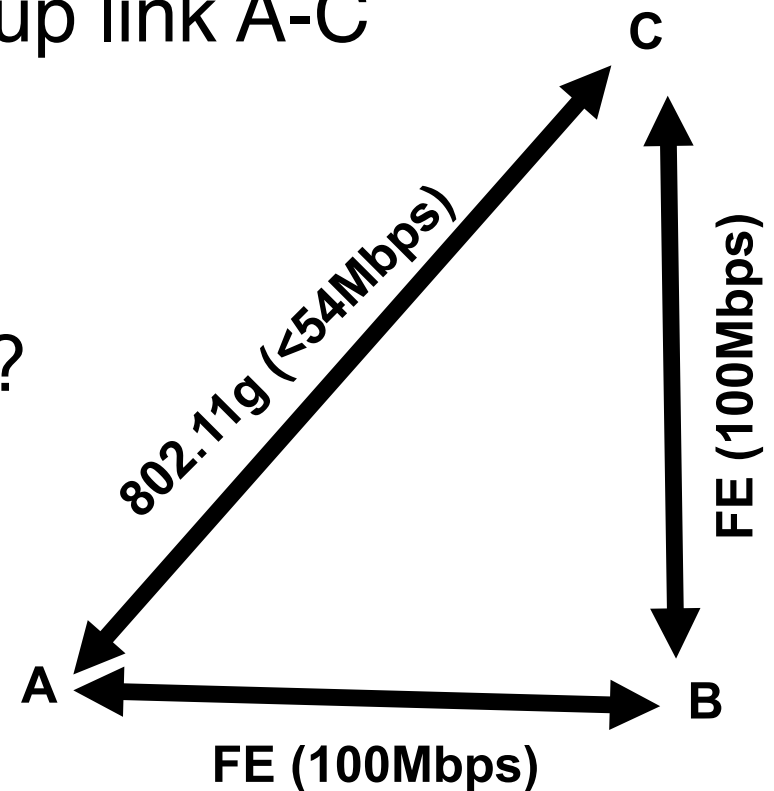
- Early, easy, and common
  - Still used today in SOHO gear
  - Set and forget; few knobs
- Straightforward implementation of Distance Vector algorithm
  - routing by rumour
  - Alternative routes are not kept
  - Metric is hop-count, and metric of 16 indicates unreachability

# RIP Metric

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Network with Fast Ethernet links between A-B and B-C, and a 802.11g backup link A-C

What is RIP's next hop for A to C?



# RIP Limitations (1)

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- No Variable Length Subnet Masks support
  - e.g. /28, but is supported in RIPv2
- Slow convergence
  - Full advertisements broadcast every 30s
  - Count-to-Infinity



- **RtrA** loses connection to **NetA**
- Before **RtrA** sends update (of break), **RtrB** sends full update, advertising **NetA** at a cost of 1
- **RtrA** now thinks it can get to **NetA** via **RtrB** at a cost of  $1+1=2$  hops
- **RtrA** advertises this to **RtrB**, which sees an increased cost. **RtrB** advertises  $2+1 = 3$ , continues to 16 (infinity/unreachability)



# RIP Limitations (2)

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- Limited diameter of network (15 hops)
- Only metric is hop-count
  - poor support for heterogeneous networks
- Advertisements not authenticated (v1)
  - attacking the network made very easy

## **RIP is okay when ...**

- Small network or minimal IT support
- Homogeneous network
  - With respect to link speed
- Convergence time is acceptable

# RIP Versions

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- RIP V1
  - No support for variable length subnet masks (VLSM)
  - No support for router authentication
- RIP V2
  - Support Classless Inter-Domain Routing (CIDR)
  - Support authentication (MD5)
- RIPng
  - RIP next generation: an extension of RIPv2 to support IPv6
  - No router authentication. IP routers were supposed to use IPsec for authentication

# RIP Optimizations

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- Split Horizon
  - prevent a router from advertising a route back onto the interface from which it was learned.
- Hold Down Timer
  - upon failure, hold the route as down until the network has converged, to expire old information
- Poison Reverse
  - Route can be deleted due to split horizon or route timeout.
  - A router actively advertises routes as unreachable over the interface over which they were learned
- Triggered Update
  - send updates as soon as something changes about a route's metric or state

# Link State vs. Distance Vector

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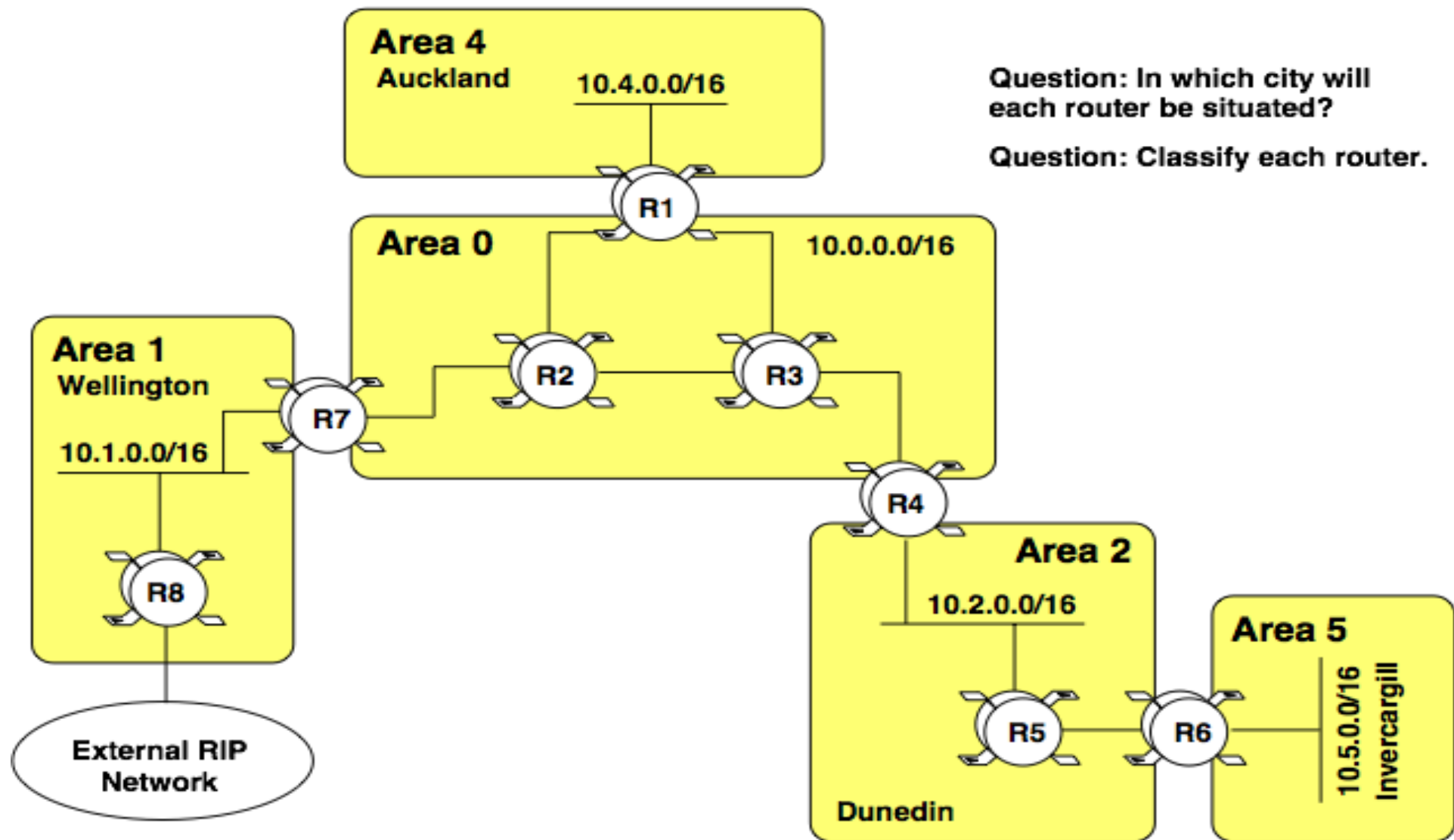
- Distance Vector
  - routing by rumour
  - unaware of topology
  - distributed Bellman-Ford algorithm
- Link State
  - multicast local link states to all routers
  - each router then assembles topology
  - Dijkstra's algorithm
  - higher memory and processor requirements

# OSPF

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- Open Shortest Path First
  - Link State protocol
- Break AS into smaller areas
  - simplifying calculation requirements
  - routes summarised crossing areas
- Areas attached to backbone
  - area 0 is backbone area

# OSPF



# Route Summarisation

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- Reduce size of routing tables
  - supernetting (route aggregation)
- Reflects hierarchical nature of a network
  - 192.168.1.32/28 + 0010 0000
  - 192.168.1.48/28 = 0011 0000
  - 192.168.1.32/27 .001

we own everything in 192.168.1.32/27

# Route Summarisation

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- Very large, complex, well-designed networks with skilled staff
- High memory and CPU requirements
- Fast convergence (but not fastest)

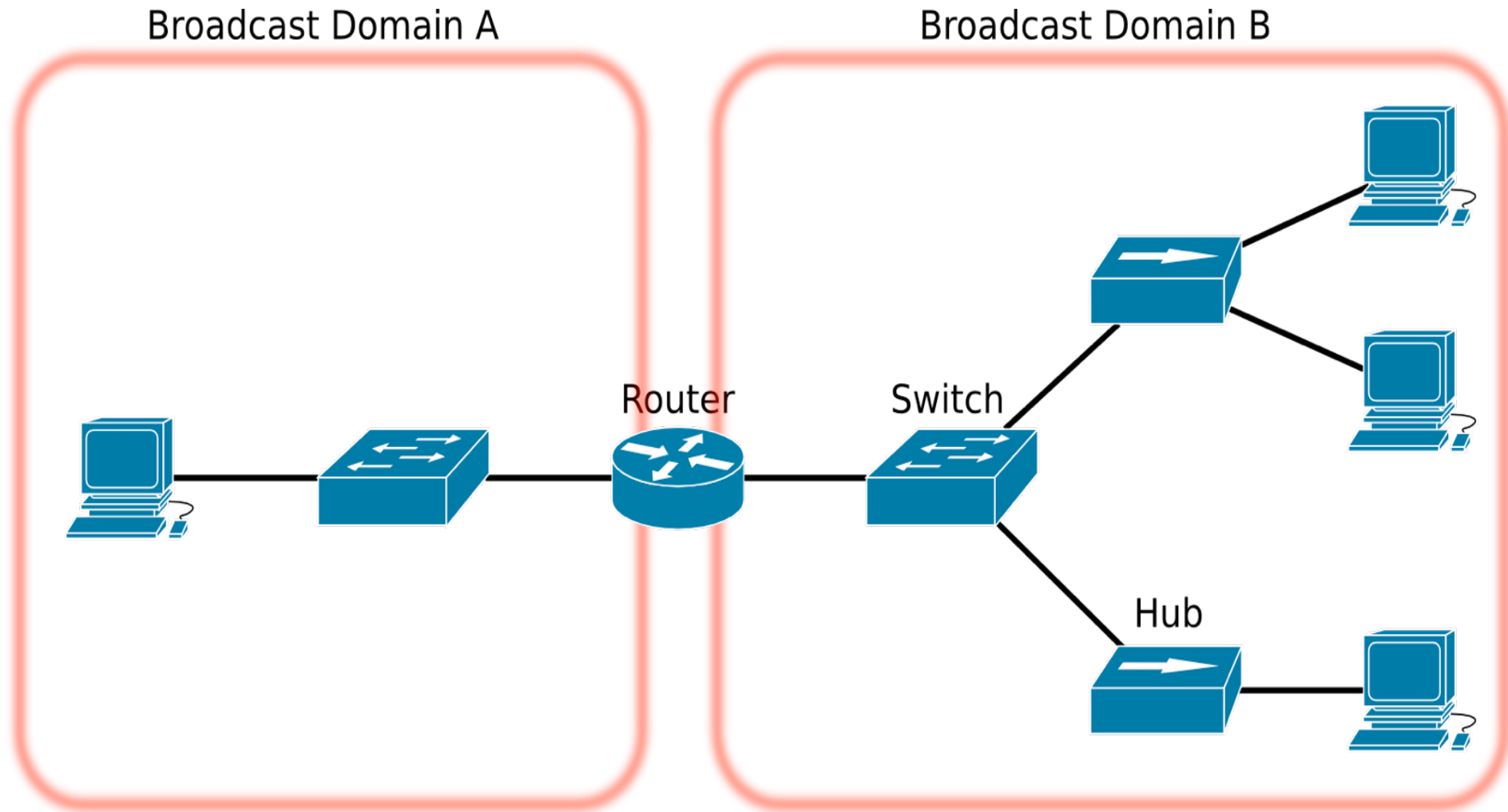


# Virtual LANs

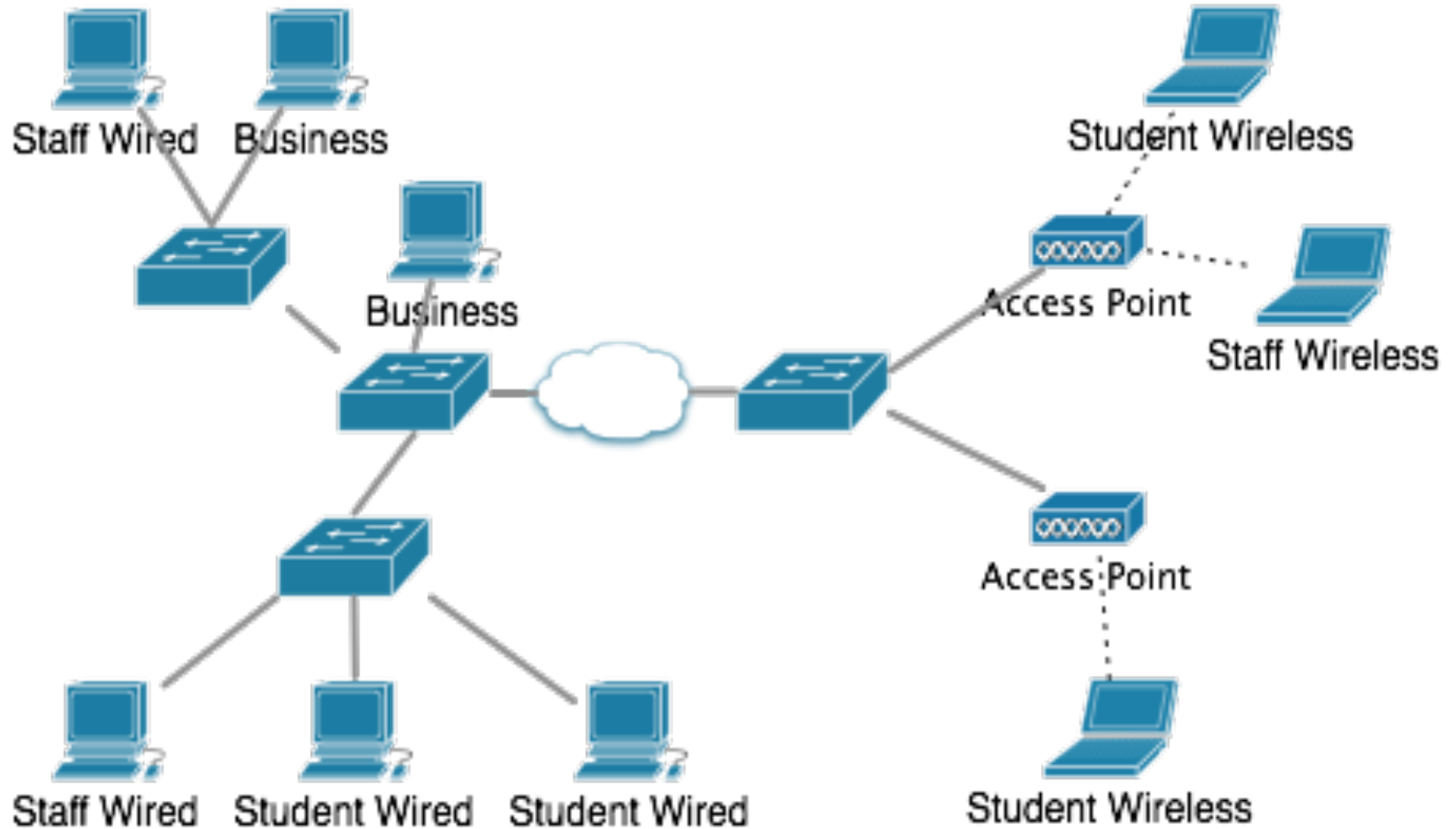
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- A local area network configured by software, not by physical wiring
  - Virtually connecting devices in different physical LANs
- IEEE802.1Q
  - The protocol most commonly used today to configure VLANs

# Broadcast Domains



# VLAN Motivations



# Virtual LAN Types

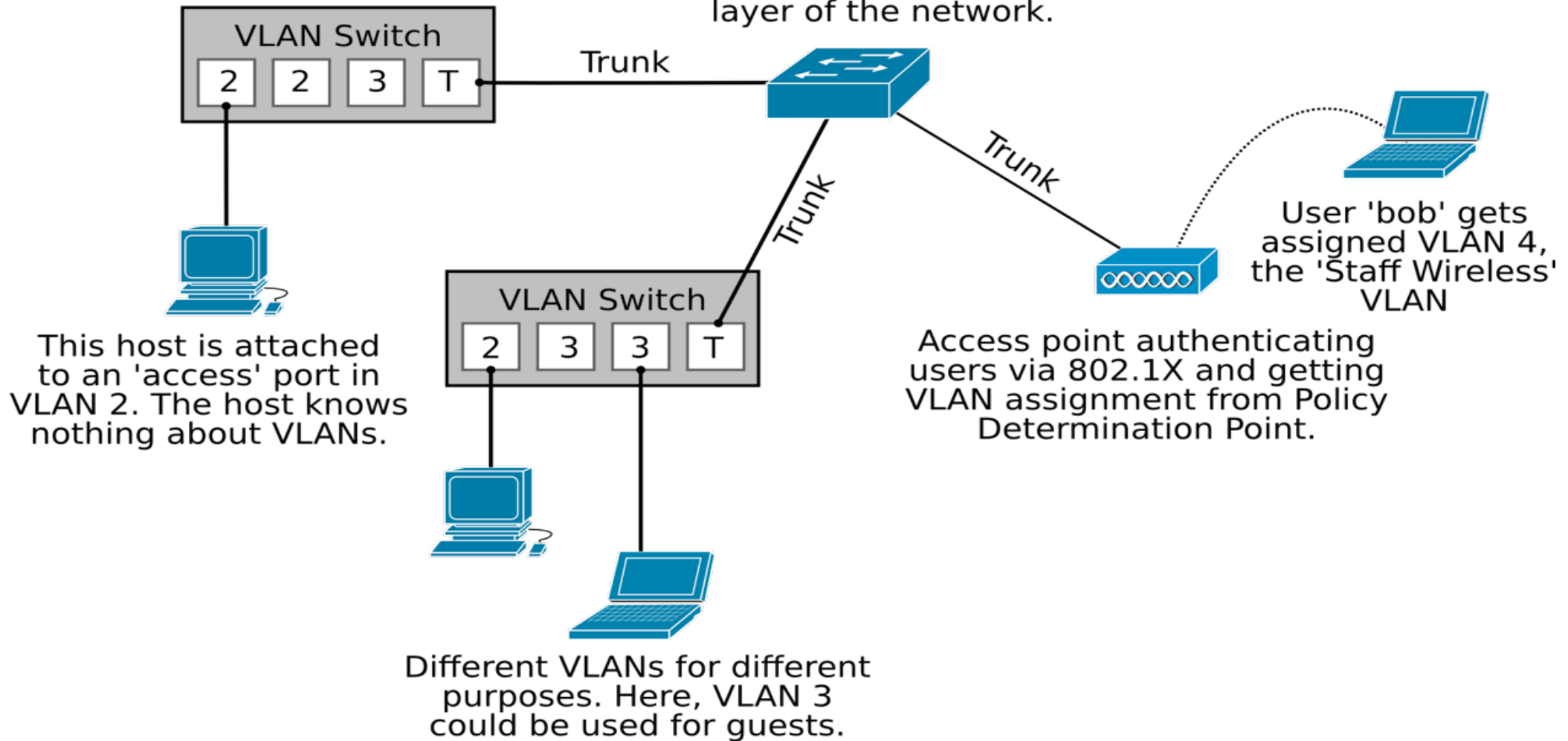
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- Layer 1 VLAN: Membership by Port
  - Does not allow for user mobility
- Layer 2 VLAN: Membership by MAC Address
  - Support user mobility
  - VLAN membership must be assigned initially.
- Layer 3 VLAN: Membership by IP Subnet Address
  - Support user mobility
  - takes longer to forward packets using Layer 3 information than using MAC addresses.

# Port Assignments

Most switch ports will be assigned manually by the network admin.

Traffic on trunking links are tagged with the appropriate VLAN. The tags are not seen by the hosts in the access layer of the network.



**You will experiment with VLANs  
and internal routing in labs**