Simple Permutations

R.L.F. Brignall

joint work with Sophie Huczynska, Nik Ruškuc and Vincent Vatter

School of Mathematics and Statistics University of St Andrews

Thursday 15th June, 2006

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

Introduction

Basic Concepts

- Permutation Classes
- Intervals and Simple Permutations
- 2 Algebraic Generating Functions for Sets of Permutations
 - Finitely Many Simples
 - Sets of Permutations
- 3 A Decomposition Theorem with Enumerative Consequences

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

- Aim
- Pin Sequences
- Decomposing Simple Permutations
- 4 Decidability and Unavoidable Structures
 - More on Pins
 - Decidability

Outline



Basic Concepts

- Permutation Classes
- Intervals and Simple Permutations
- 2 Algebraic Generating Functions for Sets of Permutations
 - Finitely Many Simples
 - Sets of Permutations
- 3 A Decomposition Theorem with Enumerative Consequences
 - Aim
 - Pin Sequences
 - Decomposing Simple Permutations
- 4 Decidability and Unavoidable Structures
 - More on Pins
 - Decidability

- Regard a permutation of length *n* as an ordering of the symbols 1,..., *n*.
- A permutation τ = t₁t₂...t_k is involved in the permutation σ = s₁s₂...s_n if there exists a subsequence s_{i1}, s_{i2},..., s_{ik} order isomorphic to τ.

Example

- Regard a permutation of length *n* as an ordering of the symbols 1,..., *n*.
- A permutation τ = t₁ t₂... t_k is involved in the permutation σ = s₁ s₂... s_n if there exists a subsequence s_{i1}, s_{i2},..., s_{ik} order isomorphic to τ.

Example

- Regard a permutation of length *n* as an ordering of the symbols 1,..., *n*.
- A permutation τ = t₁ t₂... t_k is involved in the permutation σ = s₁ s₂... s_n if there exists a subsequence s_{i1}, s_{i2},..., s_{ik} order isomorphic to τ.



- Regard a permutation of length *n* as an ordering of the symbols 1,..., *n*.
- A permutation τ = t₁ t₂... t_k is involved in the permutation σ = s₁ s₂... s_n if there exists a subsequence s_{i1}, s_{i2},..., s_{ik} order isomorphic to τ.



- Involvement forms a partial order on the set of all permutations.
- Downsets of permutations in this partial order form permutation classes.
- A permutation class C can be seen to avoid certain permutations. Write C = Av(B).

Example

The class C = Av(12) consists of all the decreasing permutations:

 $\{1, 21, 321, 4321, \ldots\}$

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

- Involvement forms a partial order on the set of all permutations.
- Downsets of permutations in this partial order form permutation classes.
- A permutation class C can be seen to avoid certain permutations. Write C = Av(B).

Example

The class C = Av(12) consists of all the decreasing permutations:

 $\{1, 21, 321, 4321, \ldots\}$

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

- Involvement forms a partial order on the set of all permutations.
- Downsets of permutations in this partial order form permutation classes.
- A permutation class C can be seen to avoid certain permutations. Write C = Av(B).

Example

The class C = Av(12) consists of all the decreasing permutations:

 $\{1, 21, 321, 4321, \ldots\}$

(日) (日) (日) (日) (日) (日) (日) (日)

- Involvement forms a partial order on the set of all permutations.
- Downsets of permutations in this partial order form permutation classes.
- A permutation class C can be seen to avoid certain permutations. Write C = Av(B).

Example

The class C = Av(12) consists of all the decreasing permutations:

 $\{1, 21, 321, 4321, \ldots\}$

(日) (日) (日) (日) (日) (日) (日) (日)

- Involvement forms a partial order on the set of all permutations.
- Downsets of permutations in this partial order form permutation classes.
- A permutation class C can be seen to avoid certain permutations. Write C = Av(B).

Example

The class C = Av(12) consists of all the decreasing permutations:

 $\{1,21,321,4321,\ldots\}$

Permutation Classes

Generating Functions

• C_n – permutations in C of length n.

• $\sum |C_n| x^n$ is the generating function.

Example

The generating function of C = Av(12) is:

$$1+x+x^2+x^3+\cdots = \frac{1}{4}$$

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ □ ● ● ●

Generating Functions

- C_n permutations in C of length n.
- $\sum |C_n|x^n$ is the generating function.

Example

The generating function of C = Av(12) is:

$$1 + x + x^2 + x^3 + \cdots = \frac{1}{4}$$

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ □ ● ● ●

Generating Functions

- C_n permutations in C of length n.
- $\sum |C_n|x^n$ is the generating function.

Example

The generating function of C = Av(12) is:

$$1+x+x^2+x^3+\cdots = \frac{1}{1}$$

▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

Generating Functions

- C_n permutations in C of length n.
- $\sum |C_n|x^n$ is the generating function.

Example

The generating function of C = Av(12) is:

$$1 + x + x^2 + x^3 + \cdots = \frac{1}{1 - x}$$

▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

• Pick any permutation π .

An interval of π is a set of contiguous indices *I* = [*a*, *b*] such that π(*I*) = {π(*i*) : *i* ∈ *I*} is also contiguous.



- Pick any permutation π .
- An interval of π is a set of contiguous indices *I* = [*a*, *b*] such that π(*I*) = {π(*i*) : *i* ∈ *I*} is also contiguous.



- Pick any permutation π .
- An interval of π is a set of contiguous indices *I* = [*a*, *b*] such that π(*I*) = {π(*i*) : *i* ∈ *I*} is also contiguous.



- Pick any permutation π .
- An interval of π is a set of contiguous indices *I* = [*a*, *b*] such that π(*I*) = {π(*i*) : *i* ∈ *I*} is also contiguous.



Simple Permutations

• Only intervals are singletons and the whole thing.

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

Example

Simple Permutations

• Only intervals are singletons and the whole thing.





▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

Simple Permutations

• Only intervals are singletons and the whole thing.

Example



▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 のへで

Simple Permutations

• Only intervals are singletons and the whole thing.





▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

Simple Permutations

• Only intervals are singletons and the whole thing.





▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

Simple Permutations

• Only intervals are singletons and the whole thing.



・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

Simple Permutations

• Only intervals are singletons and the whole thing.





▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

Intervals and Simple Permutations

Special Simple Permutations



Parallel alternations.

- Wedge permutations
- Two flavours of wedge simple permutation.

Intervals and Simple Permutations

Special Simple Permutations



ヘロン 人間と 人間と 人間と 一間

- Parallel alternations.
- Wedge permutations
- Two flavours of wedge simple permutation.

Intervals and Simple Permutations

Special Simple Permutations



・ロット (雪) (日) (日) (日)

- Parallel alternations.
- Wedge permutations not simple!
- Two flavours of wedge simple permutation.

Intervals and Simple Permutations

Special Simple Permutations



◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

- Parallel alternations.
- Wedge permutations
- Two flavours of wedge simple permutation.

Outline

Basic Concepts

- Permutation Classes
- Intervals and Simple Permutations

2 Algebraic Generating Functions for Sets of Permutations

- Finitely Many Simples
- Sets of Permutations
- 3 A Decomposition Theorem with Enumerative Consequences
 - Aim
 - Pin Sequences
 - Decomposing Simple Permutations
- 4 Decidability and Unavoidable Structures
 - More on Pins
 - Decidability

Algebraic Generating Functions for Sets of Permutations Finitely Many Simples





Algebraic Generating Functions for Sets of Permutations Finitely Many Simples





Algebraic Generating Functions for Sets of Permutations Finitely Many Simples





Motivation II

"...the standard intuition of what a family with an algebraic generating function looks like: the algebraicity suggests that it may (or should...), be possible to give a recursive description of the objects based on disjoint union of sets and concatentation of objects."

- Bousquet-Mélou, 2006

(日) (日) (日) (日) (日) (日) (日) (日)

- We can always write permutations with a simple block pattern, the substitution decomposition.
- Use recursive enumeration for classes with finitely many simple permutations.
- Expect an algebraic generating function.
Motivation II

"...the standard intuition of what a family with an algebraic generating function looks like: the algebraicity suggests that it may (or should...), be possible to give a recursive description of the objects based on disjoint union of sets and concatentation of objects."

- Bousquet-Mélou, 2006

(日) (日) (日) (日) (日) (日) (日) (日)

- We can always write permutations with a simple block pattern, the substitution decomposition.
- Use recursive enumeration for classes with finitely many simple permutations.
- Expect an algebraic generating function.

Motivation II

"...the standard intuition of what a family with an algebraic generating function looks like: the algebraicity suggests that it may (or should...), be possible to give a recursive description of the objects based on disjoint union of sets and concatentation of objects."

- Bousquet-Mélou, 2006

- We can always write permutations with a simple block pattern, the substitution decomposition.
- Use recursive enumeration for classes with finitely many simple permutations.
- Expect an algebraic generating function.

Motivation II

"...the standard intuition of what a family with an algebraic generating function looks like: the algebraicity suggests that it may (or should...), be possible to give a recursive description of the objects based on disjoint union of sets and concatentation of objects."

- Bousquet-Mélou, 2006

(日) (日) (日) (日) (日) (日) (日) (日)

- We can always write permutations with a simple block pattern, the substitution decomposition.
- Use recursive enumeration for classes with finitely many simple permutations.
- Expect an algebraic generating function.

Theorem (RB, SH, VV)

- the number of permutations in C_n (Albert and Atkinson),
- the number of alternating permutations in C_n ,
- the number of even permutations in C_n ,
- the number of **Dumont** permutations in C_n ,
- the number of permutations in C_n avoiding any finite set of blocked or barred permutations,
- the number of involutions in C_n , and
- Any (finite) combination of the above.

Theorem (RB, SH, VV)

- the number of permutations in C_n (Albert and Atkinson),
- the number of alternating permutations in C_n ,
- the number of even permutations in C_n ,
- the number of **Dumont** permutations in C_n ,
- the number of permutations in C_n avoiding any finite set of blocked or barred permutations,
- the number of involutions in C_n , and
- Any (finite) combination of the above.

Theorem (RB, SH, VV)

- the number of permutations in C_n (Albert and Atkinson),
- the number of alternating permutations in C_n ,
- the number of even permutations in C_n ,
- the number of **Dumont** permutations in C_n ,
- the number of permutations in C_n avoiding any finite set of blocked or barred permutations,
- the number of involutions in C_n, and
- Any (finite) combination of the above.

Theorem (RB, SH, VV)

- the number of permutations in C_n (Albert and Atkinson),
- the number of alternating permutations in C_n ,
- the number of even permutations in C_n ,
- the number of **Dumont** permutations in C_n ,
- the number of permutations in C_n avoiding any finite set of blocked or barred permutations,
- the number of involutions in C_n , and
- Any (finite) combination of the above.

Theorem (RB, SH, VV)

- the number of permutations in C_n (Albert and Atkinson),
- the number of alternating permutations in C_n ,
- the number of even permutations in C_n ,
- the number of Dumont permutations in C_n ,
- the number of permutations in C_n avoiding any finite set of blocked or barred permutations,
- the number of involutions in C_n, and
- Any (finite) combination of the above.

Theorem (RB, SH, VV)

- the number of permutations in C_n (Albert and Atkinson),
- the number of alternating permutations in C_n ,
- the number of even permutations in C_n ,
- the number of Dumont permutations in C_n ,
- the number of permutations in C_n avoiding any finite set of blocked or barred permutations,
- the number of involutions in C_n , and
- Any (finite) combination of the above.

Theorem (RB, SH, VV)

- the number of permutations in C_n (Albert and Atkinson),
- the number of alternating permutations in C_n ,
- the number of even permutations in C_n ,
- the number of Dumont permutations in C_n ,
- the number of permutations in C_n avoiding any finite set of blocked or barred permutations,
- the number of involutions in C_n , and
- Any (finite) combination of the above.

Outline

Basic Concepts

- Permutation Classes
- Intervals and Simple Permutations
- 2 Algebraic Generating Functions for Sets of Permutations
 - Finitely Many Simples
 - Sets of Permutations

3 A Decomposition Theorem with Enumerative Consequences

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

- Aim
- Pin Sequences
- Decomposing Simple Permutations
- 4 Decidability and Unavoidable Structures
 - More on Pins
 - Decidability

Aim

What we want to find



- Large simple permutation, size f(k).
- Find two simple permutations inside, each of size *k*.

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

Overlap of at most two points – almost disjoint.

Aim

What we want to find



- Large simple permutation, size f(k).
- Find two simple permutations inside, each of size *k*.

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

Overlap of at most two points – almost disjoint.

Aim

What we want to find



- Large simple permutation, size f(k).
- Find two simple permutations inside, each of size *k*.

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

Overlap of at most two points – almost disjoint.

Aim

What we want to find



- Large simple permutation, size f(k).
- Find two simple permutations inside, each of size *k*.

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

• Overlap of at most two points - almost disjoint.

Why I

• Every simple of length \geq 4 contains 132.

- Every simple of length ≥ f(4) contains 2 almost disjoint copies of 132.
- $\geq f(f(4))$ contains 4 copies of 132.

Theorem (Bóna; Mansour and Vainshtein)

Why I

- Every simple of length \geq 4 contains 132.
- Every simple of length ≥ f(4) contains 2 almost disjoint copies of 132.
- $\geq f(f(4))$ contains 4 copies of 132.

Theorem (Bóna; Mansour and Vainshtein)

Why I

- Every simple of length \geq 4 contains 132.
- Every simple of length ≥ f(4) contains 2 almost disjoint copies of 132.
- $\geq f(f(4))$ contains 4 copies of 132.

Theorem (Bóna; Mansour and Vainshtein)

Why I

- Every simple of length \geq 4 contains 132.
- Every simple of length ≥ f(4) contains 2 almost disjoint copies of 132.
- $\geq f(f(4))$ contains 4 copies of 132.

Theorem (Bóna; Mansour and Vainshtein)

Why II

• Av $(\beta_1^{\leq r_1}, \beta_2^{\leq r_2}, \dots, \beta_k^{\leq r_k})$ — the class with: $\leq r_1$ copies of β_1 , $\leq r_2$ copies of β_2 , etc.

Corollary

If the class $Av(\beta_1, \beta_2, ..., \beta_k)$ contains only finitely many simple permutations then for all choices of nonnegative integers $r_1, r_2, ..., and r_k$, the class $Av(\beta_1^{\leq r_1}, \beta_2^{\leq r_2}, ..., \beta_k^{\leq r_k})$ also contains only finitely many simple permutations.

Corollary

For all r and s, every subclass of $Av(2413^{\leq r}, 3142^{\leq s})$ contains only finitely many simple permutations and thus has an algebraic generating function.

• Av $(\beta_1^{\leq r_1}, \beta_2^{\leq r_2}, \dots, \beta_k^{\leq r_k})$ — the class with: $\leq r_1$ copies of β_1 , $\leq r_2$ copies of β_2 , etc.

Corollary

If the class $Av(\beta_1, \beta_2, ..., \beta_k)$ contains only finitely many simple permutations then for all choices of nonnegative integers $r_1, r_2, ..., and r_k$, the class $Av(\beta_1^{\leq r_1}, \beta_2^{\leq r_2}, ..., \beta_k^{\leq r_k})$ also contains only finitely many simple permutations.

Corollary

For all r and s, every subclass of $Av(2413^{\leq r}, 3142^{\leq s})$ contains only finitely many simple permutations and thus has an algebraic generating function.

• Av $(\beta_1^{\leq r_1}, \beta_2^{\leq r_2}, \dots, \beta_k^{\leq r_k})$ — the class with: $\leq r_1$ copies of β_1 , $\leq r_2$ copies of β_2 , etc.

Corollary

If the class $Av(\beta_1, \beta_2, ..., \beta_k)$ contains only finitely many simple permutations then for all choices of nonnegative integers $r_1, r_2, ..., and r_k$, the class $Av(\beta_1^{\leq r_1}, \beta_2^{\leq r_2}, ..., \beta_k^{\leq r_k})$ also contains only finitely many simple permutations.

Corollary

For all r and s, every subclass of $Av(2413^{\leq r}, 3142^{\leq s})$ contains only finitely many simple permutations and thus has an algebraic generating function.

Pin Sequences

Proper Pin Sequences



- Start with any two points.
- Extend up, down, left, or right this is a right pin.
- A proper pin must be maximal and cut the previous pin, but not the rectangle.

・ロット (雪) (日) (日) (日)

A right-reaching pin sequence.

Pin Sequences

Proper Pin Sequences



Start with any two points.

- Extend up, down, left, or right this is a right pin.
- A proper pin must be maximal and cut the previous pin, but not the rectangle.

・ロト・日本・日本・日本・日本・日本

• A right-reaching pin sequence.

Pin Sequences

Proper Pin Sequences



Start with any two points.

- Extend up, down, left, or right this is a right pin.
- A proper pin must be maximal and cut the previous pin, but not the rectangle.

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

A right-reaching pin sequence.

Pin Sequences

Proper Pin Sequences



- Start with any two points.
- Extend up, down, left, or right this is a right pin.
- A proper pin must be maximal and cut the previous pin, but not the rectangle.

▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

A right-reaching pin sequence.

Pin Sequences

Proper Pin Sequences



- Start with any two points.
- Extend up, down, left, or right this is a right pin.
- A proper pin must be maximal and cut the previous pin, but not the rectangle.

・ロット (雪) (日) (日) (日)

• A right-reaching pin sequence.

Pin Sequences

Proper Pin Sequences



- Start with any two points.
- Extend up, down, left, or right this is a right pin.
- A proper pin must be maximal and cut the previous pin, but not the rectangle.
- A right-reaching pin sequence.

Pin Sequences

Proper Pin Sequences



- Start with any two points.
- Extend up, down, left, or right this is a right pin.
- A proper pin must be maximal and cut the previous pin, but not the rectangle.
- A right-reaching pin sequence.

Pin Sequences

Proper Pin Sequences



- Start with any two points.
- Extend up, down, left, or right this is a right pin.
- A proper pin must be maximal and cut the previous pin, but not the rectangle.

・ロト ・雪 ト ・ ヨ ト

A right-reaching pin sequence.

Pin Sequences

Simples and Pin Sequences



• The points of the proper pin sequence form a simple permutation.

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

A Technical Theorem

Theorem

Every simple permutation of length at least $2(2048k^8)^{(2048k^8)^{(2k)}}$ contains either a proper pin sequence of length at least 2k or a parallel alternation or a wedge simple permutation of length at least 2k.

- Proper pin sequence \Rightarrow two almost disjoint simples.
- Parallel alternation \Rightarrow two almost disjoint simples.
- Wedge simple permutation \Rightarrow two almost disjoint simples.

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで

A Technical Theorem

Theorem

Every simple permutation of length at least $2(2048k^8)^{(2048k^8)^{(2k)}}$ contains either a proper pin sequence of length at least 2k or a parallel alternation or a wedge simple permutation of length at least 2k.

- Proper pin sequence \Rightarrow two almost disjoint simples.
- Parallel alternation \Rightarrow two almost disjoint simples.
- Wedge simple permutation \Rightarrow two almost disjoint simples.

(日) (日) (日) (日) (日) (日) (日) (日)

A Technical Theorem

Theorem

Every simple permutation of length at least $2(2048k^8)^{(2048k^8)^{(2k)}}$ contains either a proper pin sequence of length at least 2k or a parallel alternation or a wedge simple permutation of length at least 2k.

- Proper pin sequence \Rightarrow two almost disjoint simples.
- Parallel alternation \Rightarrow two almost disjoint simples.
- Wedge simple permutation \Rightarrow two almost disjoint simples.

(日) (日) (日) (日) (日) (日) (日) (日)

A Technical Theorem

Theorem

Every simple permutation of length at least $2(2048k^8)^{(2048k^8)^{(2k)}}$ contains either a proper pin sequence of length at least 2k or a parallel alternation or a wedge simple permutation of length at least 2k.

- Proper pin sequence \Rightarrow two almost disjoint simples.
- Parallel alternation \Rightarrow two almost disjoint simples.
- Wedge simple permutation \Rightarrow two almost disjoint simples.

The Decomposition Theorem

Theorem (RB, SH, VV)

There is a function f(k) such that every simple permutation of length at least f(k) contains two simple subsequences, each of length at least k, which share at most two entries in common.

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ のQで
Outline

Basic Concepts

- Permutation Classes
- Intervals and Simple Permutations
- 2 Algebraic Generating Functions for Sets of Permutations
 - Finitely Many Simples
 - Sets of Permutations
- 3 A Decomposition Theorem with Enumerative Consequences

- Aim
- Pin Sequences
- Decomposing Simple Permutations
- Decidability and Unavoidable Structures
 - More on Pins
 - Decidability

More on Pins

The Language of Pins



Encode as: 1

- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



Encode as: 11

- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



Encode as: 11R

- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



Encode as: 11RU

- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



Encode as: 11RUL

- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



• Encode as: 11RULD

- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



• Encode as: 11RULDR

- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



Encode as: 11RULDRU

- Pattern involvement \leftrightarrow partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



- Encode as: 11RULDRU
- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

More on Pins

The Language of Pins



- Encode as: 11RULDRU
- Pattern involvement ↔ partial order on pin words.
- Avoiding a pattern ↔ avoiding every pin word generating that pattern.

Decidability

Theorem (RB, NR, VV)

It is decidable whether a finitely based permutation class contains only finitely many simple permutations.

Proof.

- Technical theorem ⇒ only look for arbitrary parallel or wedge simple permutations, or proper pin sequences.
- Parallel and wedge simple permutations easily verified.
- Proper pin sequences ↔ the language of pins.
- Language of pins avoiding a given pattern is regular.

Decidable if a regular language is infinite.



It is decidable whether a finitely based permutation class contains only finitely many simple permutations.

- Technical theorem ⇒ only look for arbitrary parallel or wedge simple permutations, or proper pin sequences.
- Parallel and wedge simple permutations easily verified.
- Language of pins avoiding a given pattern is regular.
- Decidable if a regular language is infinite.



It is decidable whether a finitely based permutation class contains only finitely many simple permutations.

- Technical theorem ⇒ only look for arbitrary parallel or wedge simple permutations, or proper pin sequences.
- Parallel and wedge simple permutations easily verified.
- Proper pin sequences \leftrightarrow the language of pins.
- Language of pins avoiding a given pattern is regular.
- Decidable if a regular language is infinite.



It is decidable whether a finitely based permutation class contains only finitely many simple permutations.

- Technical theorem ⇒ only look for arbitrary parallel or wedge simple permutations, or proper pin sequences.
- Parallel and wedge simple permutations easily verified.
- Proper pin sequences ↔ the language of pins.
- Language of pins avoiding a given pattern is regular.
- Decidable if a regular language is infinite.



It is decidable whether a finitely based permutation class contains only finitely many simple permutations.

- Technical theorem ⇒ only look for arbitrary parallel or wedge simple permutations, or proper pin sequences.
- Parallel and wedge simple permutations easily verified.
- Proper pin sequences ↔ the language of pins.
- Language of pins avoiding a given pattern is regular.
- Decidable if a regular language is infinite.



It is decidable whether a finitely based permutation class contains only finitely many simple permutations.

- Technical theorem ⇒ only look for arbitrary parallel or wedge simple permutations, or proper pin sequences.
- Parallel and wedge simple permutations easily verified.
- Proper pin sequences ↔ the language of pins.
- Language of pins avoiding a given pattern is regular.
- Decidable if a regular language is infinite.

Áttu eitthvað ódýrara?

◆□▶ ◆□▶ ◆ □▶ ◆ □▶ ─ □ ─ の < @

Áttu eitthvað ódýrara? Do you have anything cheaper?

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ - 三 - のへぐ