

# Sorting with a Forklift

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- A *fork stack* is a stack whose push and pop operations are allowed to accept, or produce, sequences of elements.



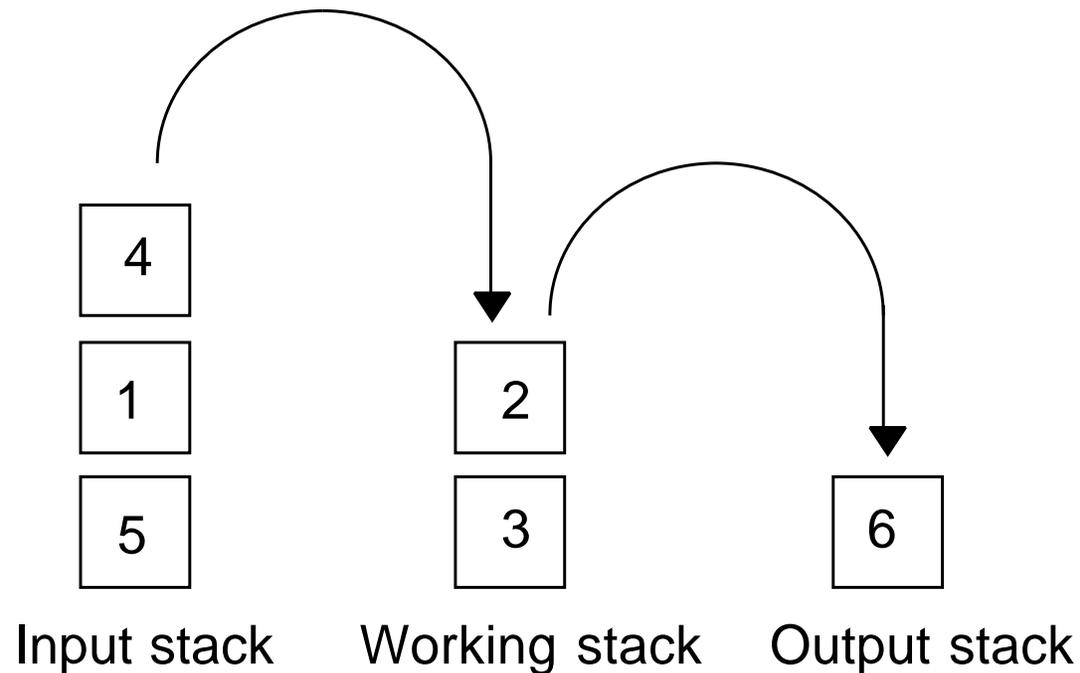
# What is a fork stack?

- Why can't we lift more than one dish from a stack?
- Why can't we put more than one dish onto a stack?
- A *fork stack* is a stack whose push and pop operations are allowed to accept, or produce, sequences of elements.
- A useful generalisation for example in merging the contents of two stacks.



# Sorting with a single forkstack

A snapshot of sorting in progress using a single forkstack (the working stack).



To complete the sort, move the pair 41 to the working stack, move 5 to the working stack and then to output, move 4 to output, and move the triple 123 to output.



# When does sorting fail?

**Definition:** For positive integers  $a$  and  $b$ ,  $a \ll b$  means  $a < b - 1$ . In a series of fork stack moves, we say that the **dreaded 13** occurs if at some point the working stack contains adjacent elements  $ab$  with  $a \ll b$ . A sequence is **near-decreasing** if it is decreasing except possibly for some steps of  $+1$ .

**Proposition:** A permutation  $\pi$  is unsortable if and only if every allowable sequence of fork stack operations that empties the input produces, at some point, the dreaded 13.



# An algorithm for sorting

**repeat**

Perform as much output as possible.

Move the maximal near-decreasing sequence from the input stack to the working stack.

**until** *input stack is empty*

**if** *working stack is empty* **then**

Success!

**else**

Failure.

**end if**



# Obstructions to sorting

- There is a finite set of patterns which, if any one of them occurs in the input stack, prevent successful sorting.
- Conversely, if none of them occur in the input stack then sorting will succeed.
- The minimal such set consists of one pattern of length 5 (35142), 45 patterns of length six, and 6 of length seven.



# Limited sorters

We can impose bounds on the sizes of either the push or pop operations (or both). An interesting special case is when the push operation is limited to a single element, and the pop operation may or may not be so limited.

We would like to find **generating functions** for the classes in this case, i.e. expressions of the form:

$$\sum_{n=0}^{\infty} \left( \begin{array}{c} \text{number of sortable} \\ \text{sequences of length } n \end{array} \right) x^n.$$



# Generating functions

Let  $f_k$  be the generating function for sortable sequences when the push operation is limited to one element, and the pop is limited to  $k$ .

In a sortable sequence,  $\pi$ , let  $t$  be the maximum for which the elements 1 through  $t$  occur in  $\pi$  in decreasing order. Then, the intervening blocks must be sortable, and the final pop of  $t$  elements must be allowed (modulo a minor quibble when 1 is the last element of  $\pi$ ).



# GF for $(1, k)$ sorting

So:

$$f_k = 1 + x f_k + (f_k - 1) \sum_{t=1}^k x^t f_k^t.$$

The different values of  $k$  provide a link from the Catalan numbers ( $k = 1$ ) which satisfy a quadratic equation, through to a solution of a different quadratic ( $k = \infty$ ).



# Growth rates

The term controlling the exponential growth of the coefficients is the reciprocal of the radius of convergence. This increases from 4 at  $k = 1$ , to 5 at  $k = \infty$ . In fact, if we write:

$$c_k = 5 - e_k$$

then

$$e_k \approx \frac{C}{3^k}.$$

The moral is, *don't pay much for extra power once you can lift six plates.*



# Other results

- If both the push and the pop operations are allowed to handle two or more objects then there are many sequences which can be sorted in more than one way.
- This complicates the enumeration problem enormously.
- For fixed bounds on the push and pop sizes we can prove: *There is a deterministic push down automaton whose accepted language contains precisely one sequence of operations to sort any sortable input sequence.*



# Late breaking news

- More sequences can be sorted using pushes and pops limited to two elements at a time, than with pushes limited to one, and pops unlimited.
- The exponential growth rate for the  $(2, 2)$  case is  $5.412\dots$ , while that for the  $(1, \infty)$  case is 5.
- The generating function  $f_{2,2}$  for the  $(2, 2)$  case satisfies an algebraic equation of degree 4, whose coefficients are polynomials in  $x$  of degree up to 4.



Thank you!