



# Simple permutations, partial well-order, and enumeration

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- Setting the scene

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- Simple permutations
- Finitely many restrictions
- Algebraic generating functions
- Some particular results

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- $A$  is *closed* if  $\sigma \in A$  and  $\pi \preceq \sigma$  implies  $\pi \in A$
- Closed sets are those of the form  $\mathcal{S}[T]$  for some  $T$

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- *Enumeration*:  $a_n = \#\pi \in A, |\pi| = n$ . Is  $a_n$  exponentially bounded? Formula for  $a_n$ ? Is  $\sum a_n x^n$  algebraic, rational?

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In many cases solutions of these problems “go together”



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A permutation with no non-trivial component is called *simple*.

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**Theorem.** *The number of simple permutations of length  $n$  is asymptotic to  $n!/e^2$ .*

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**Elementary example**

$$F = \{1, 12, 21\} \implies \mathcal{F} = \mathcal{S}[3142, 2413]$$

$$A = \mathcal{S}[231]$$

has simple set  $F$  and so  $A \subseteq \mathcal{F}$ .

# Wreath operations

Given a permutation  $\alpha$  with  $|\alpha| = n$  and permutations  $\beta_1, \dots, \beta_n$  define

$$\alpha(\beta_1, \dots, \beta_n)$$

to be the permutation  $\beta'_1 \cdots \beta'_n$  which has components  $\beta'_i$  order isomorphic to  $\beta_i$  and where the pattern of the components is  $\alpha$ .



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## Example

$$\alpha = 231, \beta_1 = 21, \beta_2 = 132, \beta_3 = 321$$

$$\alpha(\beta_1, \beta_2, \beta_3) = 54 \mid 687 \mid 321$$

# Partial well-order

**Lemma.** *If  $F$  and  $\mathcal{F}$  are as above then, for all  $\alpha \in F$  and  $\beta_1, \dots, \beta_n \in \mathcal{F}$ , we have  $\alpha(\beta_1, \dots, \beta_n) \in \mathcal{F}$ .*

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**Theorem.** *If  $F$  is finite then  $\mathcal{F}$  is partially well-ordered by pattern containment.*

**Proof**  $\mathcal{F}$  is an algebra under the wreath operations. Use Higman's theorem.

# A finite number of restrictions

**Corollary.** *Every closed class with just a finite set of simple permutations is defined by a finite number of restrictions.*

**Proof** Main step: prove that  $\mathcal{F}$  itself is defined by a finite number of restrictions. Then appeal to partial well-order.

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**Lemma.** *One of the following holds:*

- $\sigma = \alpha\beta$  where either  $\alpha < \beta$  or  $\alpha > \beta$
- $\sigma = \xi_1 \cdots \xi_m$  where each  $\xi_i$  is a maximal component and the pattern formed by the  $\xi_i$  is simple.

# Structure

Let  $F = \{1, 12, 21, \phi_3, \phi_4, \dots\}$ .

Every permutation in  $\mathcal{F} \setminus \{1\}$  has one of the forms

$$\phi(\beta_1, \dots, \beta_m)$$

where  $\phi \in F \setminus \{1\}$ .

The representation is unique except for when  $\phi = 12$  or  $\phi = 21$ .

**Notation**  $\phi(B_1, \dots, B_n)$  the set of all  $\phi(\beta_1, \dots, \beta_n)$  with  $\beta_i \in B_i$

# Structure Theorem

$F, \mathcal{F}$  as above

$\mathcal{F}_+$  (resp.  $\mathcal{F}_-$ ) permutations in  $\mathcal{F}$  with no  $\alpha\beta$  decomposition where  $\alpha < \beta$  (resp.  $\alpha > \beta$ ).

- $\mathcal{F} = \{1\} \cup 12(\mathcal{F}_+, \mathcal{F}) \cup 21(\mathcal{F}_-, \mathcal{F}) \cup \mathcal{G}$

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# Enumeration of $\mathcal{F}$

Equations satisfied by ordinary generating functions

- $f(x) = x + f_+(x)f(x) + f_-(x)f(x) + g(x)$
- $f_+(x) = x + f_-(x)f(x) + g(x)$
- $f_-(x) = x + f_+(x)f(x) + g(x)$
- $g(x) = \sum_{\phi \in F, |\phi| > 2} f(x)^{|\phi|}$

# Generating function of $\mathcal{F}$

$f(x)$  satisfies

$$f^2 - f + (x + g)(f + 1) = 0$$

In particular,

**Theorem.** *If  $|F|$  is finite,  $f(x)$  (the generating function of  $\mathcal{F}$ ) is algebraic.*

# Subclasses of $\mathcal{F}$ (for finite $F$ )

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## General idea

- Subclasses have the form  $\mathcal{F}[T]$  for some finite number of restrictions  $T$
- $\mathcal{F}$  is a disjoint union of terms of the general type  $\phi(A_1, A_2, \dots, A_n)$  where  $\phi$  is simple.

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- $\mathcal{F}$  is a disjoint union of terms of the general type  $\phi(A_1, A_2, \dots, A_n)$  where  $\phi$  is simple.

So, how do we enumerate  $\phi(A_1, \dots, A_n)[T]$ ?

# Enumeration of $\phi(A_1, \dots, A_n)[T]$

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- $\tau \preceq \alpha_1 \cdots \alpha_n$  if and only if  $\tau$  has a block decomposition  $\tau_1 \cdots \tau_l$  where each  $\tau_i \preceq \alpha_{s(i)}$  for some  $s(i)$  and the pattern of the  $\tau_i$  is a subpermutation of  $\phi$ .

# Technical calculations omitted

- $\phi(A_1, \dots, A_n)[T]$  can be expressed as a union and intersection of sets of the form  $\phi(A_1[T_1], \dots, A_n[T_n])$  where the restrictions  $T_j$  are subpermutations of the restrictions  $T$ .

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- $\phi(B_1, \dots) \cap \phi(C_1, \dots) = \phi(B_1 \cap C_1, \dots, B_n \cap C_n)$
- Use inclusion-exclusion and argue by induction

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**Theorem.** *Every proper closed subclass of  $S[231]$  has a rational generating function.*

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**Theorem.** *If the generating function of  $F$  is algebraic then the generating function of  $\mathcal{F}$  is algebraic.*